


Tensor Decompositions & Tensor Rank.

Goal: Then for order-3 tensors important primitive for studying, also higher order tensors.

Definition of Tensor Rank.

Rank-one tensors (elementary bricks)

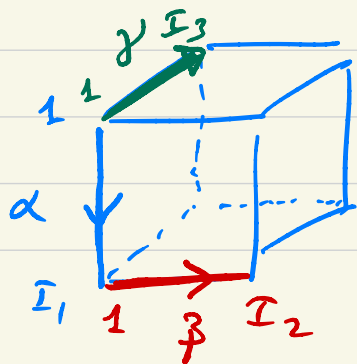
- Take three vectors $\underline{a} \in \mathbb{R}^{I_1}$, $\underline{b} \in \mathbb{R}^{I_2}$, $\underline{c} \in \mathbb{R}^{I_3}$
 a^α , $\alpha = 1 \dots I_1$, b^β and c^γ
 $\beta = 1 \dots I_2$ $\gamma = 1 \dots I_3$.

- Form the tensor product (or outer product)

$$\underline{a} \otimes \underline{b} \otimes \underline{c} = \underline{T}$$

$$T^{\alpha\beta\gamma} = a^\alpha b^\beta c^\gamma$$

3-dim array of I_1, I_2, I_3 numbers.



Rank-R tensor . $R \geq 1$.

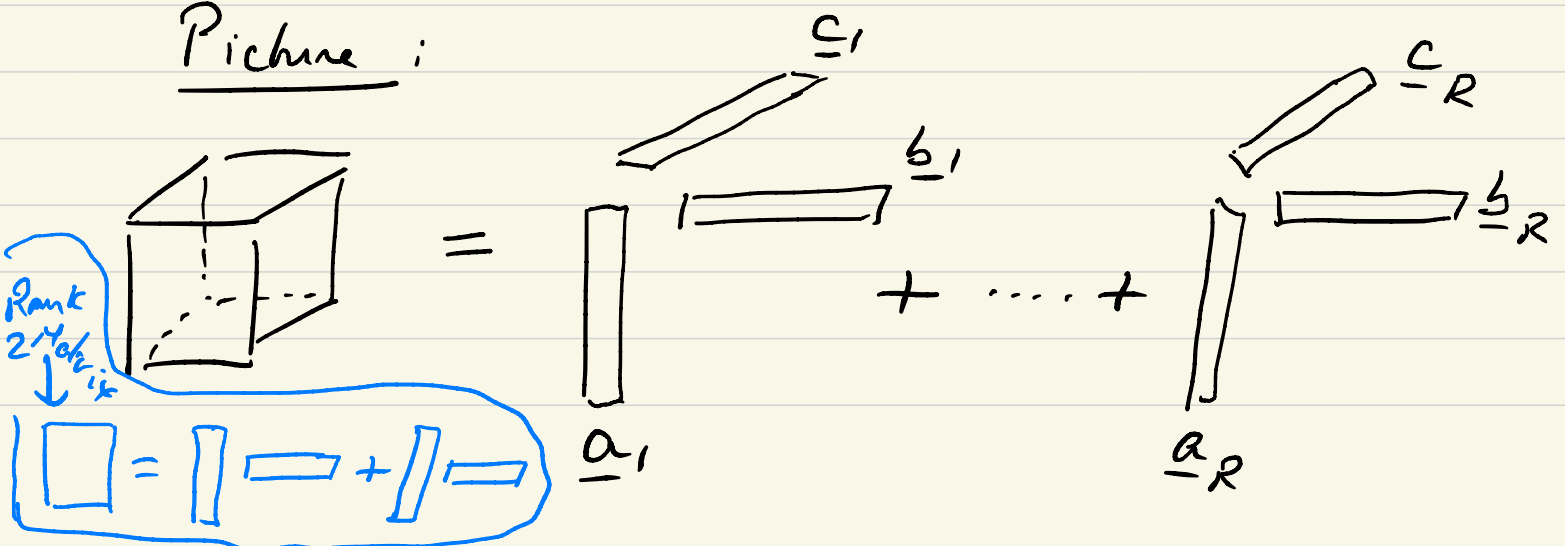
is a multi-array (3-arrays) that can be written as a sum of R rank-one tensors with R being the minimal possible number of terms.

$$T = \sum_{i=1}^{\textcircled{R}} \underline{a}_i \otimes \underline{b}_i \otimes \underline{c}_i$$

R should be the smallest possible to qualify for Tensor-Rank.

Terminology: such decomp are also called "tensor factorizations" or "polyadic decompositions".

Picture:



Adopt the following Notation:

$$T = \sum_{i=1}^R \underline{a}_i \otimes \underline{b}_i \otimes \underline{c}_i$$

$$A = [\underline{a}_1 \dots \underline{a}_R] \leftarrow I_1 \times R.$$

↑
column vectors

$$\underline{a}_1 = \begin{bmatrix} a_1^\alpha \end{bmatrix}_{\alpha=1}^{\alpha=I_1}$$

$$B = [\underline{b}_1 \dots \underline{b}_R] \leftarrow I_2 \times R$$

$$C = [\underline{c}_1 \dots \underline{c}_R] \leftarrow I_3 \times R.$$

Remarks about the notion of tensor-rank:

- 1) For ^{order-2 tensors} matrices "tensor-rank" $\Leftrightarrow$ usual lin-algebra notion of rank.

$$\dim(\text{row span}) \overset{\uparrow}{=} \dim(\text{column span}).$$

For tensors "Tensor-Rank" is NOT a linear-alg
(of order $p \geq 3$) notion.

Moreover there are other notions of Rank (see later
lects).

- 2) Non trivial to decompose a tensor in a minimal # of rank one terms.

- unicity?
- algorithms?

Rule of Thumb: R is "small" w.r.t to I_1, I_2, I_3
 $\rightarrow$ decomp tends to be unique
 $\rightarrow$ efficient algs are known.

R is "moderate"
 $\rightarrow$ decomp is unique ...
 $\rightarrow$ eff algs are not always known.
 R is "large" $\rightarrow$ don't even have unicity.

Basic Theorem. [about order 3-tensors. 1970's
Jennrich, Carroll, Harshman ...]

- Let $A = [\underline{a}_1, \dots, \underline{a}_R] \quad I_1 \times R$

$$B = [\underline{b}_1, \dots, \underline{b}_R] \quad I_2 \times R.$$

$$C = [\underline{c}_1, \dots, \underline{c}_R] \quad I_3 \times R.$$

- A & B have full column rank R . (in particular $R \leq \min(I_1, I_2)$)
- C has pairwise indep columns: $\forall i, j \quad \underline{c}_i \neq \lambda_{ij} \underline{c}_j$

- Take a 3-tensor

$$T = \sum_{i=1}^R \underline{a}_i \otimes \underline{b}_i \otimes \underline{c}_i \quad \leadsto \quad T^{\alpha\beta\gamma}$$

- Then the tensor (or multiarray of numbers $T^{\alpha\beta\gamma}$) can be decomposed in a unique way in a sum of at most R rank-one terms. And the tensor rank of T is R .

- Qualification for unicity: up to trivial rescaling of $\underline{a}, \underline{b}, \underline{c} \in [\lambda \underline{a}_i \otimes \frac{1}{\lambda} \underline{b}_i \otimes \frac{1}{\lambda} \underline{c}_i]$ & up to permutation of terms in sum.

Remarks:

- 1) The thm will be proved in a constructive way and the proof will thus give also an algorithm $\text{poly}(I_1, I_2, I_3)$.
- 2) For this thm we have $R \leq \min(I_1, I_2)$
- in this sense R is "small" -
- $\uparrow$
- uniquity & algorithmic result -

3) For $R > \min(I_1, I_2)$ less is known.

As an aside state the following result:

- suppose that K_A, K_B, K_C are Kruskal ranks of A, B, C .

(K_A is the KR if it is the max integer such that all subsets of K_A vcts in $[a_1, \dots, a_R]$ are lin-indep; $\exists K_A + 1$ subset that is lin-dep)

if $2R + 2 \leq K_A + K_B + K_C$ in decomp of T
Then this decomp is unique.

- For example $I_1 = I_2 = I_3 = N$ & $R > N$
imagine $K_A = K_B = K_C = N \Rightarrow$

here you have uniquity result for
"Moderate rank" $N < R < \frac{3N}{2} - 2$

BUT NO EFFICIENT ALGO'S ARE KNOWN!

$$R < \frac{3N - 2}{2}$$

4) This is "surprising" in the sense that for 2-tensors (metrics) it does not hold!

Recall the "relation problem":

$$A = [\underline{a}_1, \underline{a}_2] \quad \mathbb{I}_1 \times 2 \quad R = 2.$$

$$B = [\underline{b}_1, \underline{b}_2] \quad \mathbb{I}_2 \times 2 \quad R = 2.$$

A & B have rank $R = 2$.

$$T = \underline{a}_1 \otimes \underline{b}_1 + \underline{a}_2 \otimes \underline{b}_2 \quad \checkmark$$

This decomp is not unique! Why?

$$T = \underline{a}_1 \underline{b}_1^T + \underline{a}_2 \underline{b}_2^T \quad \underline{a} \otimes \underline{b} = \underline{a} \underline{b}^T$$

$$= [\underline{a}_1, \underline{a}_2] \begin{bmatrix} \underline{b}_1^T \\ \underline{b}_2^T \end{bmatrix}$$

outer product
of vectors.

$$= A B^T$$

$$= A \underbrace{R R^T}_{\mathbb{I}} B^T = \underbrace{(A R)}_A \underbrace{(B R)^T}_{B'^T}$$

BREAK NEXT
VIDEO PROVE THE
THM

$$= \underline{a}'_1 \otimes \underline{b}'_1 + \underline{a}'_2 \otimes \underline{b}'_2 \quad \checkmark \quad \#$$